FIFO and LIFO sing the BLUes*

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Abstract

FIFO, First-In-First-Out, and LIFO, Last-In-First-Out, are well-known techniques for handling sequences. In TEX macro writing they are abundant but are not easily recognized as such. TEX templates for FIFO and LIFO are given and their use illustrated. The relation with Knuth's `\dolist`, answer ex11.5, and `\ctest`, p.376, is given.

Keywords

FIFO, LIFO, list processing, plain TEX, education, macro writing.

Introduction

It started with the programming of the Tower of Hanoi in TEX, van der Laan (1992a). For printing each tower the general FIFO — First-In-First-Out[1] — approach was considered.[2] In literature (and courseware) the programming of these kind of things is done differently by each author, inhibiting intelligibility. In pursuit of Wirth (1976), TEX templates for the FIFO (and LIFO) paradigm will hopefully improve the situation.

In this article we will see various slightly different implementations of the basic FIFO principle.

FIFO

In the sequel, I will restrict the meaning of FIFO to an input stream which is processed argument-wise. FIFO can be programmed in TEX as template

```
\def\fifo#1{\ifx\ofif#1\ofif\fi
  \process#1\fifo}
\def\ofif#1\fifo{\fi}
```

The `\fifo` command calls a macro `\process` that handles the individual arguments. Often you can copy `\fifo` straight out of this article, but you have to write a version of `\process` that is specific to your application.

To get the flavor.

Length of string. An alternative to Knuth's macro `\getlength` (*The TEXbook*, p.219) is obtained via the use of `\fifo` with

```
\newcount\length
\def\process#1{\advance\length1 }
```

Then `\fifo aap  noot\ofif \number\length` yields the length 7.[3]

Number of asterisks. An alternative to Knuth's `\atest` (*The TEXbook*, p.375), for determining the number of asterisks, is obtained via `\fifo` with

```
\def\process#1{\if*#1\advance\acnt by1
  \fi}
\newcount\acnt
```

Then `\fifo abc*de*\ofif \number\acnt` yields the number of asterisks: 2.[4]

Vertical printing. David Salomon treats the problem of vertical printing in his courseware. Via an appropriate definition of `\process` and a suitable invocation of `\fifo` it is easily obtained.

```
\def\process#1{\hbox{#1}}
\vbox{\offinterlineskip\fifo abc\ofif}
```

yields a
b
c.

Tower of Hanoi. Printing of a tower can be done via

```
\def\process#1{\hbox to3ex{%
  \hss\vrule width#1ex height1ex\hss}}
\vbox{\baselineskip1.1ex\fifo12\ofif}
```

Termination. For the termination of the tail recursion the same TEXnique as given in *The TEXbook*, p.379, in the macro `\deleterightmost`, is used. This is elaborated as `\break` in Fine (1992), in relation to termination of the loop. The idea is that when `\ofif` is encountered in the input stream, that is, when `\ifx\ofif#1...` is true, all tokens in the macro up to and including `\fifo` — the start for the next level of recursion — are gobbled by a subsequent call to `\ofif`.[5] Because the matching `\fi` is gobbled too, this token is inserted via the replacement text of `\ofif`. This TEXnique is better than Kabelschacht's (1987), where the token

* Earlier versions appeared in MAPS92.1 and proceedings EuroTEX '92. BLU is Ben Lee User of *TEXbook* fame. It makes the title sing, I hope.

[1] See Knuth (1968), section 2.2.1.

[2] In the Tower of Hanoi article Knuth's list datastructure was finally used — *The TEXbook*, Appendix D.2 — with FIFO inherent.

[3] Insert `\obeyspaces` when the spaces should be counted as well.

[4] As the reader should realize, this works correctly when there are first level asterisks *only*. For counting at all levels automatically, a more general approach is needed, see Knuth's `\ctest`, p.376.

[5] In contrast with usual programming of recursion start with the infinite loop, and then insert the `\if...\ofif\fi`.

preceding the `\fi` is expanded after the `\fi` via the use of `\expandafter`. When this is applied the exchange occurs at each level in the recursion. It is also better than the `\let\nxt=...` TEXnique, which is used in *The TEXbook*, for example in `\iterate`, p.219, because there are no assignments.

My first version had the two tokens after `\ifx` reversed — a cow flew by — and made me realize the non-commutativity of the *first level* arguments of TEX's conditionals. For example, `\ifx aa\empty...` differs from `\ifx\empty aa...`, and `\if\ab\aa...` from `\if\aa\ab...`, with `\def\aa{aa}`, `\def\ab{ab}`. In math, and in programming languages like Pascal, the equality relation is commutative,[6] and no such thing as expansion comes in between. When you are not alert with respect to expansion, TEX's `\if`-s can surprise you.

The `\fifo` macro is a basic one. It allows one to proceed along a list — at least conceptually — and to apply a (user) specified process to each list element. By this approach the programming of going through a list is *separated* from the various processes to be applied to the elements.[7] It adheres to the *separation of concerns* principle, which I consider fundamental.

The input stream is processed argument-wise, with the consequence that first level braces will be gobbled. Beware! Furthermore, no outer control sequences are allowed, nor `\par`-s. The latter can be permitted via the use of `\long\def`.

A general approach — relieved from the restrictions on the input stream: *every token* is processed until `\ofif` — is given in *The TEXbook*, answer ex11.5 (`\dolist...`) and on p.376 (`\ctest...`). After adaptation to the `\fifo` notation and to the use of macros instead of token variables, Knuth's `\dolist` comes down to

```
\def\fifo{\afterassignment\tap
  \let\nxt= }
\def\tap{\ifx\nxt\ofif\ofif\fi\process
  \nxt\fifo}
\def\ofif#1\fifo{\fi}
```

This general approach is indispensable for macro writers. My less general approach can do a lot already, for particular applications, as will be shown below. But, ... beware of its limitations.

Variations. The above `\fifo` can be seen as a template for encoding tail recursion in TEX, with arguments taken from the input stream one after another. An extension is to take two arguments from the input stream at a time, with the second argument to look ahead, via

```
\def\fifo#1#2{\process#1\ifx\ofif#2
  \ofif\fi\fifo#2}
\def\ofif#1\ofif{\fi}
```

Note the systematics in the use of the parameter separator in `\ofif`; here `\ofif` and in the previous macro `\fifo`, the last token of the replacement text. Although the principle of looking ahead with recursion is abundant in computer programming, a small example to illustrate its use is borrowed from Salomon: delete last character of argument. It is related to `\deleterightmost` (*The TEXbook*p.379). Effective is the following, where a second parameter for `\fifo` is introduced to look ahead, which is inserted back when starting the next recursion level

```
\def\gobblelast#1{\fifo#1\ofif}
\def\fifo#1#2{\ifx\ofif#2\ofif\fi%
  #1\fifo#2}
\def\ofif#1\ofif{\fi}
```

Then `\gobblelast{aap}` will yield `aa`.
And what about recursion without parameters? A nice example of that is a variant implementation of Knuth's `\iterate` of the `\loop` (*The TEXbook*, p.219).

```
\def\iterate{\body\else\etareti\fi%
  \iterate}
\def\etareti#1\iterate{\fi}
```

This `\iterate` contains only 5 tokens in contrast with Knuth's 11. The efficiency and the needed memory is determined by the number of tokens in `\body`, and therefore this 5 vs. 11 is not relevant. The idea behind including this variant here is that the FIFO principle can lead to simple encoding of tail recursion even when no arguments are processed.

Variable number of parameters. TEX macros can take at most 9 parameters. The above `\fifo` macro can be seen as a macro which is relieved from that restriction. Every group, or admissible token, in the input stream after `\fifo` up to and including `\ofif`, will become an argument to the macro.

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[6] So are TEX's `\if`-s after expansion.

[7] If a list has to be *created*, Knuth's list datastructure might be used, however, simplifying the execution of the list. See *The TEXbook*, Appendix D.2.

When the `\ofif` token is reached, the recursion— that is reading of arguments—will be terminated.[8]

Unknown number of arguments. Tutelaers (1992), as mentioned by Eijkhout (1991), faced the problem of inputting a chess position. The problem is characterized by an unspecified number of positions of pieces, with for the pawn positions the identification of the pawn generally omitted. Let us denote the pieces by the capital letters K(ing), Q(ueen), B(ishop), (k)N(ight), R(ook), and P(awn), with the latter symbol default. The position on the board is indicated by a letter a, b, c, ..., or h, followed by a number, 1, 2, ..., or 8. Then, for example,

```
\position{Ke1, Qd1, Na1, e2, e4}
```

should entail the invocations

```
\piece{K}{e1}\piece{Q}{d1}\piece{N}{a1}
\piece{P}{e2}\piece{P}{e4}
```

This can be done by an appropriate definition of `\position`, and an adaptation of the `\fifo` template, via

```
\def\position#1{\fifo#1,\ofif,}
\def\fifo#1,{\ifx\ofif#1\ofif
  \fi\process#1\relax\fifo}
\def\ofif#1\fifo{\fi}
\def\process#1#2#3{\ifx\relax#3
  \piece{P}{#1#2}\else\piece#1{#2#3}\fi}
```

With the following definition (simplified in relation to Tutelaers)

```
\def\piece#1#2{ #1-#2}
```

we get K-e1 Q-d1 N-a1 P-e2 P-e4.

For an unknown number of arguments at two levels see the Nested FIFO section.

Citation lists. In a list of citations it is a good habit to typeset three or more consecutive numbers as a range. For example 1, 2, 3 as 1–3. This must be done via macros when the numbers are represented by symbolic names, which get their value on the fly. In general the sequence must be sorted[9] before typesetting. This has been elaborated by Arseneau (1992) in a few LaTeX styles, and for plain TeX by myself. I used the FIFO paradigm in the trivial, stepping-stone, variant of typesetting an explicit non-descending sequence in range notation. The resulting 'process'-macro could be used in the general case, once I realized that FISO—First-In-Smallest-Out—was logically related to FIFO: the *required* elements are yielded one after the other, whether the first, the last, the smallest, or ... you name it. Perhaps this is a nice exercise for the reader. For a solution see van der Laan (1993).

Vowels, voilà. Schwarz (1987) coined the problem to print vowels in bold face.[10] The problem can be split into two parts. First, the general part of going character by character through a string, and second, decide whether the character at hand is a vowel or not.

For the first part use `\fifo` (or Knuth's `\dolist`).

For the second part, combine the vowels into a string, `aeiou`, and the problem can be reduced to the question $\langle char \rangle \in$ `aeiou`? Earlier, I used this approach in searching a card in a bridge hand, van der Laan (1990, the macro `\strip`). That was well-hidden under several piles of cards, I presume? The following encoding is related to `\ismember` (*The TeXbook*, p.379).

```
\newif\iffound
\def\loc#1#2{%locate #1 in #2
  \def\locate##1#1##2\end{\ifx\empty##2%
    \empty\foundfalse\else\foundtrue\fi}%
  \locate#2#1\end}
```

Then `\fifo Audacious\ofif` yields **A**ud**a**c**iou**s, with

```
\def\process#1{\uppercase{\loc#1}%
  {AEIOU}\iffound{\bf#1}\else#1\fi}
```

Variation. If in the invocation `\locate#2#1` a free symbol is inserted between `#2` and `#1`, then `\loc` can be used to locate substrings.[11] And because $\{string_1 \in string_2\} \wedge \{string_2 \in string_1\} \Rightarrow string_1 = string_2$, the variant can be used for the equality test for strings. See also the Multiple FIFO subsection, for general and more effective alternatives for equality tests of strings.

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[8] Another way to circumvent the 9 parameters limitation is to associate names to the quantities to be used as arguments, let us say via def's, and to use these quantities via their names in the macro. This is Knuth's parameter mechanism and is functionally related to the so-called keyword parameter mechanism of command languages, and for example ADA.

[9] The sorting of short sequences within TeX has been elaborated by Jeffreys 1990, and myself in Syntactic Sugar.

[10] His solution mixes up the picking up of list elements and the process to be applied. Moreover, his nesting of `\if`-s in order to determine whether a character is a vowel or not, is not elegant. Fine (1992)'s solution, via a switch, is not elegant either.

[11] Think of finding 'bb' in 'ab' for example, which goes wrong without the extra symbol.

Processing lines. What about processing lines of text? In official, judicial, documents it is a habit to fill out lines of text with dots.[12] This can be solved by making the end-of-line character active, with the function to fill up the line. A general approach where we can `\process` the line, and not only append to it, can be based upon `\fifo`.

One can wonder, whether the purpose can't be better attained, while using TEX as formatter, by filling up the last line of paragraphs by dots, because TEX justifies with paragraphs as units.

Processing words. What about handling a list of words? This can be achieved by modifying the `\fifo` template into a version which picks up words, `\fifow`, and to give `\processw` an appropriate function.

```
\def\fifow#1 {\ifx\wofif#1\wofif\fi
  \processw{#1}\ \fifow}
\def\wofif#1\fifow{\fi}
```

Underlining words . In print it is uncommon to emphasize words by underlining. Generally another font is used; see discussion of exercise 18.26 in *The TEXbook*. However, now and then people ask for (poor man's) underlining of words. The following `\processw` definition underlines words picked up by `\fifow`

```
\def\processw#1{\vtop{\hbox{\strut#1}
  \hrule}}
```

Then

```
\leavevmode\fifow leentje leerde lotje
lopen langs de lange lindenlaan \wofif
\unskip.
```

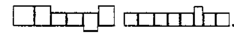
yields leentje leerde lotje lopen langs de lange lindenlaan.

Nested FIFO

One can nest the FIFO paradigm. For processing lines word by word, or words character by character.

Words character by character. Ex11.5 can be solved by processing words character by character. A solution to a slightly simplified version of the exercise reads

```
\fifow Though  exercise \wofif \unskip.
%with
\def\processw#1{\fifo#1\ofif}
\def\process#1{\boxit#1}
\def\boxit#1{\setbox0=\hbox{#1}\hbox
  {\lower\dp0\vbox{\offinterlineskip\hrule
  \hbox{\vrule\phantom#1\vrule}\hrule}}}
```

yields .

In the spirit of `\dolist...`, ex11.5, is

```
%variant neglecting word structure
\def\fifo{\afterassignment\tap
  \let\nxt= }
\def\tap{\ifx\nxt\ofif\ofif
  \fi\process\nxt\fifo}
\def\ofif#1\fifo{\fi}
\def\process#1{\if\space\nxt\
  \else\boxit#1\fi}
\fifo Though    exercise\ofif.
```

with the same result .

Mark up natural data. Data for `\h(v)align` needs `&` and `\cr` marks. We can get plain TEX to append a `\cr` at each (natural) input line (*The TEXbook*, p.249). An extension of this is to get plain TEX to insert `\cs`-s, column separators, and `\rs`-s, row separators, and eventually to add `\lr`, last row, at the end, in natural data. For example prior to an invocation of `\halign`, one wants to get plain TEX to do the transformation

P*ON
DEK* ⇒ P\cs*\csO\csN\rsD\csE\csK\cs*\lr

This can be done via

```
$$\vcenter{\hbox{P*ON}\kern.5ex
  \hbox{DEK*}} \,\Rightarrow\,
%And now right, mark up part
\bdata P*ON
DEK*
\edata\markup\data
\vcenter{\hbox{\data}}$$
```

with

```
\def\bdata{\bgroup\obeylines\store}
\def\store#1\edata{\egroup\def\data{#1}}
\def\markup#1{\ea\xdef\ea#1\ea{\ea
              \fifol#1\lofif}}
```

and auxiliaries

```
\let\nx=\noexpand
{\catcode`\^^M=13
\gdef\fifol#1^^M#2{\fifo#1\ofif%
  \ifx\lofif#2\nx\lr\lofif
  \fi\nx\rs\fifol#2}}
\def\lofif#1\lofif{\fi}
\def\fifo#1#2{#1\ifx\ofif#2\ofif
  \fi\nx\cs\fifo#2}
\def\ofif#1\ofif{\fi}
%with for this example
\def\cs{{\sevenrm{\tt\char92}cs}}
\def\rs{{\sevenrm{\tt\char92}rs}}
```

[12] The problem was posed at EuroTEX '91 by Theo Jurriens.

```
\def\lr{{\sevenrm{\tt\char92}lr}}
```

The above came to mind when typesetting crosswords,[13] van der Laan (1992b,-d),[14] while striving after the possibility to allow natural input, independent of `\halign` processing.

Multiple FIFO

What about FIFO for more than one stream?[15] For example comparing strings, either for equality or with respect to lexicographic ordering? Eijkhout (1992, p.137, 138) provided for these applications the macros
`\ifAllChars...\Are...\TheSame`,
and
`\ifallchars...\are...\bfore`.
The encodings are focused at mouth processing. The latter contains many `\expandafter`-s.

A basic approach is: loop through the strings character by character, and compare the characters until either the assumed condition is no longer true, or the end of either one of the strings, has been reached.

Equality of strings. The TEX-specific encoding, where use has been made of the property of `\ifx` for control sequences, reads

```
\def\eq#1#2{\def\st{#1}\def\nd{#2}
\ifx\st\nd\eqtrue\else\eqfalse\fi}
```

with auxiliary `\newif\ifeq`.

As a stepping stone for lexicographic comparison, consider the general encoding

```
\def\eq#1#2{\continuetrue\eqtrue
  \loop\ifx#1\empty\continuefalse\fi
       \ifx#2\empty\continuefalse\fi
  \ifcontinue \nxte#1\nxtt \nxte#2\nxtu
       \ifx\nxtt\nxtu
       \else\eqfalse\continuefalse\fi
  \repeat
  \ifx\empty#1\ifx\empty#2
  \else\eqfalse\fi\else\eqfalse\fi}
```

with auxiliaries

```
\newif\ifcontinue\newif\ifeq
\def\nxte#1#2{\def\pop##1##2\pop{%
  \gdef#1{##2}\gdef#2{##1}}\ea\pop#1\pop}
```

Then

```
\def\t{abc}\def\u{ab}
\eq\t\u\ifeq$abc=ab$\else$abc\not=ab$\fi
```

yields $abc \neq ab$.

Lexicographic comparison. Assume that we deal with lower case and upper case letters only. The encoding of `\sle`—String Less or Equal—follows the same flow as the equality test, `\eq`, but differs in the test, because of TEX's expansion mechanisms

```
\def\sle#1#2{%#1, #2 are def's
  \global\sletrue\global\eqtrue
  {\continuetrue
  \loop\ifx#1\empty\continuefalse\fi
       \ifx#2\empty\continuefalse\fi
  \ifcontinue\nxte#1\nxtt\nxte#2\nxtu
       \ea\ea\ea\lle\ea\nxtt\nxtu
  \repeat}
  \ifeq\ifx\empty#2\ifx\empty#1
        \else\global\slefalse\fi\fi
  \fi}
```

with auxiliaries (`\lle`=Letter Less or Equal)

```
\newif\ifcontinue
\global\newif\ifsle\global\newif\ifeq
\def\nxte#1#2{\def\pop##1##2\pop{%
  \xdef#1{##2}\xdef#2{##1}}\ea\pop#1\pop}
\def\lle#1#2{\uppercase{\ifnum`#1=`#2}
  \else\continuefalse\global\eqfalse
    \uppercase{\ifnum`#1>`#2}{}\global
                          \slefalse\fi
  \fi}
```

For example

```
\def\t{ABC}\def\u{ab}\sle\t\u
\ifsle$ABC\le ab$\else$ABC>ab$\fi
```

yields $ABC > ab$;

```
\def\t{aa}\def\u{a}\sle\t\u
\ifsle$aa\le a$\else$aa>a$\fi
```

yields $aa > a$;

```
\def\t{aa}\def\u{b}\sle\t\u
\ifsle$aa\le b$\else$aa>b$\fi
```

yields $aa \leq b$;

```
\def\t{noo}\def\u{apen}\sle\t\u
\ifsle$noo⟨apen£ “else£noo⟩apen$\fi
```

yields $noo > apen$.

The above can be elaborated with respect to `\read` for strings each on a separate file, to strings with accented letters, to the inclusion of an ordering table, and in general to sorting. Some of the mentioned items will be treated in Sorting in BLUe, to come.

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13 With *, or ␣, given an appropriate function.

14 In the latter article I set the puzzles via direct use of nested FIFO. No `\halign` use nor mark up phase.

15 For simplicity the streams are stored in def-s, because `\read` inputs lines.

LIFO

A modification of the `\fifo` macro — `\process{#1}` invoked at the end instead of at the beginning — will yield the Last-In-First-Out template. Of course LIFO can be applied to reversion on the fly, without explicitly allocating auxiliary storage.[16]

```
\def\lifo#1#2\ofil{\ifx\empty#2
  \empty\ofil\fi\lifo#2\ofil\process#1}
\def\ofil#1\ofil{\fi}
```

The test for emptyness of the second argument is similar to the TEXnique used by Knuth in `\displaytest` (*The TEXbook*, p.376): `\if!#3!...`.

With the identity — `\def\process#1{#1}`, or the invoke `\process#1` replaced by `#1`[17] — the template can be used for reversion on the fly For example `\lifo aap\ofil` yields `paa`.[18]

Change of radix. In *The TEXbook* a LIFO exercise is provided at p.219: print the digits of a number in radix 16 representation. The encoding is based upon the property

$$(N \div r^k) \bmod r = d_k, \quad k = 0, 1, \ldots, n,$$

with radix r, coefficients d_k, and the number representation

$$N = \sum_{k=0}^{n} d_k\, r^k.$$

There are two ways of generating the numbers d_k: starting with d_n, or the simpler one starting with d_0, with the disadvantage that the numbers are generated in reverse order with respect to printing. The latter approach is given in *The TEXbook*, p.219. Adaptation of the LIFO template does not provide a solution much different from Knuth's, because the numbers to be typeset are generated in the recursion and not available in the input stream.

[16] Johannes Braams drew my attention to Knuth and MacKay (1987), which contained among others `\reflect...\tcelfer`. They compare `#1` with `\empty`, which is nice. The invocation needs an extra token, `\empty` — a so-called sentinel, see Wirth (1976) — to be included before `\tcelfer`, however. (Knuth and Mackay hide this by another macro which invokes `\reflect...\empty\tcelfer`). My approach requires at least one argument, with the consequence that the empty case must be treated separately, or a sentinel must be appended after all.

[17] Remember the stack size limitations.

[18] Note that Knuth's test `\if!#3!...` goes wrong for `#3` equals `!`, and similarly my use of the idea goes wrong for `#2` equals `\empty`, which is not 'empty'. Given the context those situations don't occur, however.

Acknowledgements

Włodek Bzyl and Nelson Beebe are kindly acknowledged for their help in clearing up the contents and correcting my use of English, respectively.

Conclusion

In looking for a fundamental approach to process elements sequentially — not to confuse with list processing where the list is also built up, see *The TEXbook*, Appendix D.2, or with processing of *every* token in the input stream, see ex11.5 or p.376 — TEX templates for FIFO and LIFO emerged.

The templates can be used for processing lines, words or characters. Also processing of words line by line, or characters word by word, can be handled via nested use of the FIFO principle.

The FIFO principle along with the look ahead mechanism is applied to molding natural data into representations required by subsequent TEX processing.

Courseware might benefit from the FIFO approach to unify answers of the exercises of the macro chapter.

TEX's `\ifx...` and `\if...` conditionals are non-commutative with respect to their *first level* operands, while the similar mathematical operations are, as are the operations in current high-level programming languages.

Multiple FIFO, by comparing strings lexicographically, has been touched upon.

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LATEX

An Update on the babel System

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Abstract

This article describes the changes that have been made to the babel system since the article describing the system appeared in *TUGboat* 12, no. 2. This article announces the release of a new version of the babel system.

1 Introduction

Since the publication of the babel system in *TUGboat* [1] several changes have occurred. With the new release of LATEX — which appeared at the end of 1991 — the internationalised version ILATEX, prepared by Joachim Schrod [2], was withdrawn. But some of its functionality was still needed, so a modification of the babel system was necessary.

Besides this a couple of bugs were reported and had to be fixed. The major problem was that the language changing commands were not 'local', they contained global definitions. In the current version these commands obey grouping correctly.

Some macros that formerly were in language-specific files have been moved to the core of the system, because they are being used in several language-specific files.

2 Changes to the core of babel

The changes to the core of the babel system are the most extensive.

`\selectlanguage`

The babel user-command `\selectlanguage` now also accepts a control sequence as its argument. This was included to provide compatibility for users who were used to the syntax of the original `german.tex`, but wanted to switch to babel. The escape character is 'peeled off' and the name of the control sequence is then used as the name of the language to select.

Another change to the `\selectlanguage` macro is that it now stores the name of the current language in the control sequence `\languagename`. The