

# Symbolic Computation for Electronic Publishing

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## Abstract

Recently, we developed and reported some novel ways to make MATHEMATICA produce derivations that contain conventionally structured narratives and formulas, using a procedure that interprets files containing expressions, which MATHEMATICA evaluates symbolically, control information and text. Now, this work is extended to give  $\TeX$  typeset output. It supports the interactive crafting of publications that contain mechanically generated formulas and diagrams, for teaching and research, in subjects that use mathematics and other algorithmic methods. It goes beyond the direct capabilities of the built-in `Splice` and `TeXForm` functions.

## 1. Introduction

Recent dramatic gains in the accessibility and power of workstations has given symbolic computation greatly increased exposure. Several monographs describe MATHEMATICA (Wolfram 1991) and other computer algebra systems. Interactive use, with and without “notebooks”, and batch runs that produce lengthy FORTRAN statements for numerical evaluation on mainframe computers are common. A relatively underutilized application is the combined use of electronic typesetting and symbolic computation to produce research journals, monographs, text books, reference compendia, problem books and other documents containing mathematical formulas that are:

- the result of lengthy proofs and derivations,
- burdensome to copy and check,
- numerous and closely related to each other,
- needed in different notations,
- internally repetitive (e.g. matrices whose elements are written in a common form).

Computerized symbolic calculation of the formulas in all these cases can save considerable time and effort that is needed, otherwise, to produce the results, type the manuscripts, and read and correct the proofs. The need to print mathematical formulas that were produced by primitive symbolic computations led to some of the early work on electronic typesetting over 30 years ago (Barnett 1965)—the formulas were too numerous and lengthy to derive by

hand. The present paper describes some current activity that uses MATHEMATICA with plain  $\TeX$  (Knuth 1986),  $\LaTeX$  (Lamport 1994) and  $\mathcal{A}\mathcal{M}\mathcal{S}\text{-}\TeX$  (Spivak 1986). It builds on our `autorun` procedure and `bi10` package (Barnett 1991; Barnett 1993; Barnett 1994), and the application of these to atomic and molecular structure calculations (Barnett 1991a; Barnett 1994a), robot kinematics (Chen 1991) and other topics.

To introduce some of the principles that underly this work, the box below depicts a workstation screen during a simple MATHEMATICA session. The default prompt for input has the form “`In[n]:=`”. The value that the system computes for each input statement is displayed alongside an “`Out[n]=`” tag. Here, statement 1 assigns a simple expression to `s`. Statement 2 replaces `n` by 1 and expands the result.

```
In[1]:= s = (1 + x/(x+1))^(n+1)

Out[1]= (1 +  $\frac{x}{1+x}$ )1+n

In[2]:= s /. n -> 1 // Expand

Out[2]= 1 +  $\frac{x^2}{(1+x)^2}$  +  $\frac{2x}{1+x}$ 
```

MATHEMATICA evaluates:

$$a /. b \rightarrow c$$

by substituting  $c$  for  $b$  throughout  $a$ . Also,  $b \rightarrow c$  is called a “transformation rule”, a “replacement rule” or just a “rule”. The function  $f$  of  $x$  is written as:

$$f[x], \quad x//f, \quad f@x$$

These are synonymous.  $f$  is called their “head”.

## 2. Interpreting a control file

Our typesetting work builds on the `autorun` procedure (Barnett 1991), that reads a file of MATHEMATICA statements and, in the default mode, mimics interactive operation. Thus, given the control file `demo1` consisting of the two records:

```
s = (1 + x/(x+1))^(n+1)
s /. n -> 1 // Expand
```

a continuation of the MATHEMATICA session of Section 1 is shown in the next box. Only the `autorun` statement is typed. It simulates the earlier interactive action, by interpreting the control file and generating “In” and “Out” tags that combine the sequence numbers in the MATHEMATICA and `autorun` sessions. The screen display can be recorded, e.g., as a UNIX script file, and printed.

```
In[3]:= autorun[demo1]
In[3.1]:= s = (1 + x/(x+1))^(n+1)
Out[3.1]= 
$$\left(1 + \frac{x}{1+x}\right)^{n+1}$$

In[3.2]:= s /. n -> 1 // Expand
Out[3.2]= 
$$1 + \frac{x^2}{(1+x)^2} + \frac{2x}{1+x}$$

```

Often, users want to see the output of the successive symbolic evaluations, without a playback of the input. `autorun` provides an “outputOnly” mode, and lets the control file contain formatting and related “directives” that do not get displayed, and text that does. These are labeled by # and \* symbols, as in the file `demo2`, shown in the next box. It contains two `bilo` formatting functions. The first reverses the reordering of the exponent — `numbersLast` simply moves numerical terms to the right in its target. The second treats the target as a polynomial in  $x$ .

```
# outputOnly
* Given
s = (1 + x/(x+1))^(n+1)
# format = toThe[n+1][numbersLast]
* then the expansion, when n=1, is
# format = sortToIncreasePowersOf[x]
s /. n -> 1 // Expand
```

Interpreting this file gives:

```
In[4]:= autorun[demo2]
Given

$$\left(1 + \frac{x}{1+x}\right)^{n+1}$$

then the expansion, when n=1, is

$$1 + \frac{2x}{1+x} + \frac{x^2}{(1+x)^2}$$

```

## 3. Producing typeset output

A file `demo2.tex` was produced from `demo2` by:

```
autorecord[demo2]
```

This invokes `autorun` in a mode that converts the evaluated expressions to  $\text{\TeX}$ . `autorecord` then writes the relevant portion of the screen display to an output file, invokes  $\text{\LaTeX}$  or, optionally,  $\text{\TeX}$ , and a previewer. Then it prints, if requested. In the  $\text{\TeX}$  file that produced this paper, `\input demo2` imported the  $\text{\TeX}$  coded contents of the next box.

```
Given

$$\left(1 + \frac{x}{1+x}\right)^{n+1}$$

then the expansion, when n=1, is

$$1 + \frac{2x}{1+x} + \frac{x^2}{(1+x)^2}$$

```

The “ $n=1$ ” in the second line of text is set in math mode by putting delimiters around it in the input.

A file close to `demo2.tex` can be produced from a slightly different control file using the built-in `Splice` function of MATHEMATICA, which uses the built-in `TeXForm` function to do the  $\text{\TeX}$  encoding of individual expressions. We use the procedure `toTeX`, that is part of our `forTeX` package, to do the encoding. It provides greater flexibility and control.

### 4. Compound heads

Formatting the final result in demo2 involved, implicitly, the evaluation of:

```
sortToIncreasePowersOf[x][1 + ...]
```

In this, the head itself is a function of x.

```
sortToIncreasePowersOf[x]
```

We call this a "compound head". In general:

```
f[u1, u2, ...][v1, v2, ...]
```

denotes a function  $g[v_1, v_2, \dots]$ , where  $g$  is the function  $f[u_1, u_2, \dots]$ . Thus, the entire expression is a function of  $u_1, u_2, \dots, v_1, v_2, \dots$ . The principle is extended, e.g. in `a[b][c][d][e]`. Invented by Ruffini two centuries ago (see (Ruffini 1799) and (Cajori 1919, page 81)), the notation was reinvented by Curry (Curry et al. 1958; Hindley and Seldin 1986). We use it extensively, for convenience both in symbolic computation and when encoding to T<sub>E</sub>X. The ability of MATHEMATICA to handle it is very important.

### 5. The basic toTeX conventions

`autorecord` hides the T<sub>E</sub>X output of `toTeX` from the user. To show it here, we display a MATHEMATICA expression  $s$ , the "leads to" symbol  $\rightsquigarrow$  and the typeset value of `toTeX[s]` as, for example, in:

```
sin[theta] sin[pi+theta] ~ sin theta sin(pi + theta)
```

This example, incidentally, illustrates selective parenthesization which T<sub>E</sub>XForm does not provide.

For processing by `toTeX`, elementary expressions are represented using standard MATHEMATICA conventions, that include `==` between the sides of an equation. The characters listed in (Knuth 1986, Appendix F, Tables 1-14) and many other objects are denoted by T<sub>E</sub>X-like names that omit the backslash. Thus:

```
alpha, beta, ... ~ alpha, beta, ...
aleph, hbar, ... ~ N, h, ...
pm, mp, ... ~ +/-, +/-, ...
leq, prec, ... ~ <=, <, ...
```

We use function notation to represent fonts, subscripts, superscripts, decorations, special bracketing and binary expressions. Thus:

```
bf[A + B + C] ~ A + B + C
P[sub[n]][sup[m]][x] ~ P_n^m(x)
hat[x], overline[x+y] ~ x-hat, x-bar+y-bar
enpr[x], encr[x, y] ~ (x), {x, y}
cap[cap[A, B], C] ~ A union B intersection C
```

`enpr` abbreviates "enclose in parentheses" and `sapr` provides parentheses that are sized automatically. Corresponding names with `br`, `cr`, `lr` and `so` forth provide square brackets, curly braces, `<` `>` and other enclosing symbols.

We use compound heads to represent several functionals. These include:

```
sum[i, j, k][f[i]] ~ sum_{i=j}^k f(i)
limit[t -> 0][h[t]] ~ lim_{t->0} h(t)
D$[{x, 2}, y][phi] ~ d^3 phi / (dx^2 dy)
```

`bilo` has many functions that operate on these representations in symbolic calculations, e.g. `rightExpand` performs the reduction:

$$\sum_{i=j}^k f(i) \rightarrow \sum_{i=j}^{k-1} f(i) + f(k)$$

Denoting MATHEMATICA evaluation by  $\Rightarrow$ , the action of `rightExpand` is specified also by:

```
sum[i, j, k][f[i]] // rightExpand
=> sum[i, j, k-1][f[i]] + f[k]
```

Other `bilo` functions perform the reductions:

$$\sum_{i=j}^k f(i) \rightarrow \sum_{i=j}^{k+1} f(i) - f(k+1)$$

$$\sum_{i=j}^k (a + b + \dots) \rightarrow \sum_{i=j}^k a + \sum_{i=j}^k b + \dots$$

$$\sum_{i=j}^k f(i) + \sum_{i=k+1}^l f(i) \rightarrow \sum_{i=j}^l f(i)$$

and several related operations.

The compound head notation allows numerous variations, e.g. in symbolic calculations that involve integrals, we use expressions that include:

```
integral[x][y] ~ int y dx
integral[{x1, x2}][y] ~ double int y dx1 dx2
integral[{x1, x2, x3}][y] ~ triple int y dx1 dx2 dx3
integral[x, u, v][y] ~ int_u^v y dx
integral[{x1, u1, v1}, {x2, u2, v2}][y] ~ int_{u1}^{v1} int_{u2}^{v2} y dx1 dx2
integral[{x1, u1, v1}, ..., {xn, un, vn}][y] ~ int_{u1}^{v1} ... int_{un}^{vn} y dx1 ... dxn
lineIntegral[dot, L, s][bf[V]] ~ int_L V . ds
surfaceIntegral[S][f] ~ double int_S f dS
```

### 6. Operating on toTeX input

Other names can be converted to T<sub>E</sub>X-like names by replacement rules that precede the statements which

invoke `toTeX`. The rules are input as directives (see Section 9). So are targeting functions that put function heads onto specific subexpressions as in:

```
f[a] + g[b] + h[c] //
  inSuccession[
    toThe[f][bf]
    collectivelyToTerms[
      containingAny[g, h]][sabr]]
⇒ bf[f[a]] + sabr[g[b] + h[c]]
≈ f(a) + [g(b) + h(c)]
```

The targeting function here is read “apply the function `bf` to the part of the target that has the head `f`, then apply the function `sabr` collectively to the parts that contain either `g` or `h`.” `biTo` contains an easily extensible set of functions that address terms, factors, elements of a list or a vector or a matrix, coefficients, arguments, parts of a fraction or a relationship, patterns and explicit items.

**Suppressing parentheses:** `toTeX` converts the brackets in an arbitrary MATHEMATICA function to parentheses. `unpr` elides them. Thus:

```
f[x], f[unpr[x]], f[unpr[circ]][unpr[x]]
≈ f(x), fx, f ∘ x
```

**Varying the style:** Our default style for square roots is controlled by the `toTeX` variable `defaultSqrtStyle`. Initially it is 1, with the effect:

```
sqrt[1 - sqrt[1 - Delta^2]]
≈ √(1 - √(1 - Δ²))
```

Setting it to 2 leads to output in the style:

```
sqrt[1 - sqrt[1 - Delta^2]]
≈ (1 - (1 - Δ²)½)½
```

The default can be changed by a simple MATHEMATICA assignment. Also, in the formation of the `TEX` codes by `toTeX` each `sqrt[x]` is converted to the intermediate form:

```
style[sqrt, m][x]
```

with `m` set to the value of `defaultSqrtStyle` that is current. An individual square root in a formula can be set in a different style `n` by changing it to `style[sqrt, n][x]` before applying `toTeX`. To do this, the transformation rule

```
sqrt -> style[sqrt, n]
```

is applied to the appropriate subexpression(s) by a `biTo` targeting function. Similar tactics are used to style other fractional powers and fractions.

**Independent options:** Several features in the styling of an integral can be altered independently. `toTeX` begins the encoding by converting the primary head `integral` to the form:

```
style[integral, a1 -> b1, a2 -> b2, ...]
```

The system initializes the option values `bi` that are associated with the option names `ai`. The statement:

```
setOptions[integral, aj -> bj, ak -> bk, ...]
```

is used to change options that, for integrals, include:

- `differentialsInNumerator` — defaults to `no`,
- `allowRepeatedIntegral` — no outputs multiple indefinite integrals with a single  $\int$  symbol,
- `variableInLower` — `all`, `multiple`, and `none` include the variable of integration in the lower limit(s) of, respectively, all definite integrals, only multiple definite integrals, and in none,
- `variableInUpper` — `yes` and `no` include and omit the integration variable in the upper limit,
- `includeLimits` — `yes` and `no` allow and preclude the inclusion of limits.

Options are set analogously for other functionals.

## 7. `biTo` formatting functions

**Rearranging an expression:** Often, the default order of an expression imposed by MATHEMATICA does not give the preferred form. Non-atomic objects are represented internally by nested function expressions, for example,  $(a+b)(c+d)$  by:

```
Times[Plus[a, b], Plus[c, d]]
```

Formatting often involves rearranging the argument lists of one or more such functions, and insulating the result from automatic reordering. `biTo` includes a suite of sorting functions based on the action of:

```
{x1, ..., xn} // partByCriteria[c1, ..., cm]
```

This evaluates to  $\{s_1, \dots, s_{m+1}\}$  where:

1. for  $v = 1$  to  $m$ ,  $s_v$  is the list of  $x_i$  for which  $c_1(x), \dots, c_{v-1}(x)$  are false and  $c_v(x)$  is true,
2.  $s_{m+1}$  is the list of  $x_i$  for which  $c_1(x), \dots, c_m(x)$  are false.

The order of the  $x$ s in each list  $s_v$  is the same as in the original list  $\{x_1, \dots, x_n\}$ .

The higher level `sortByPresence[v1, v2, ...]` and `sortByAbsence[v1, v2, ...]` are very useful. The  $v$ 's can be explicit terminal subexpressions or intermediate level heads or patterns for either. Related functions include those used in Section 2.

The `biTo` sort functions are applied to the relevant subexpressions by targeting functions. For this kind of work, expressions containing `biTo` sort functions are much shorter, usually, than is possible with functions that require pairwise ordering criteria.

Further rearrangement requirements include moving the denominator of a fraction to the beginning, as in  $\frac{1}{2}(x+y)$  and preventing, say,  $-(a+b)$  from opening up to  $-a-b$ . `biTo` handles these.

**Skeletalizing:** The reduction:

```
(1+x)10 // Expand // showTerms[{1, 6, -1}]
⇒ 1+«4 terms»+252x5+«4 terms»+x10
```

typifies the action of another suite of `biTo` functions. If  $p$  stand for a pointer list, then `showTerms[p]`

and `showFactors[p]` act on `Plus` and `Times` expressions, `showElements[p]` acts on lists, vectors, matrices and tensors. `showArguments[p]` acts on any function. Each `showObjects` function skeletalizes the target. By default, each shows that items are omitted by an expression of the form:

«*n object(s)*»

This can be overridden to display ellipses.

**Insertion of codes:** If  $p$  is a MATHEMATICA pattern that matches one of the terms in a `Plus`, then applying `splitAfter` to this `Plus` breaks the line after the + or - sign that follows the term. `splitBefore[p]` breaks the line before the sign.

Applying a `split` function to any non-atomic expression puts the subexpression which is matched onto a new line. `insertBefore[p][c]` is wrapped by `splitBefore[p]`, where  $c$  is a string of TeX codes. `splitAfter` is handled correspondingly. Further to TeX wrappers will allow convenient control of indentation and vertical spacing. Nested targeting expressions can be used to focus on the part of the target into which codes must be inserted.

## 8. A simple forTeX application

This relates to teaching special functions of mathematical physics. The statement:

`autorecord[legendreAuto]`

reads the file `legendreAuto`, shown in the next box, and constructs the file `legendreAuto.tex`.

```
* {\bf Demonstrating orthogonality.\ } We
* integrate the product of two Legendre
* polynomials of different degree. Consider
s1 = integral[x, -1, 1][P[2, x] P[4, x]]
* Substituting explicit polynomials for
* the $P_n(x)$ gives
# beginLeftAlignedGroup["="]
s2 = s1 // evaluateAndHoldSpecialFunctions
s3 = s2 // releaseAndExpandAll
# endLeftAlignedGroup
* Then term by term integration gives
# beginRunonGroup["="]
s4 = s3 // integrateTermByTermAndHold
s4 // allowEvaluation
# endRunonGroup
```

The statement `\input legendreAuto` in the file that set this paper produced the contents of the box in the right hand column. By default, the `outputOnly` mode is in effect. So is the sandwiching of individual TeX coded output expressions between centering delimiters. The first three records in the control file begin the text. The next record is the MATHEMATICA

statement that assigns, to the variable `s1`, the integral to be evaluated. The next two records provide the next piece of text. As regards the rest of the file:

1. `beginLeftAlignedGroup` puts the codes to begin a left-aligned multi-formula display before the next coded expression, and an = symbol at the right. Also, it sets a switch that includes codes to continue the display, when subsequent input expressions are processed.
2. `evaluateAndHoldSpecialFunctions` is a short MATHEMATICA procedure, written for this application, that converts Legendre, Laguerre and other special polynomials, written as  $P[n, x]$ ,  $L[n, x]$ , ..., to explicit polynomials that are kept separate. Here, this separation stops MATHEMATICA multiplying the denominators.
3. `endLeftAlignedGroup` writes the codes to end a multi-formula display and resets switches.
4. `releaseAndExpandAll` removes the insulation around individual parts of the target expression and applies the built-in `ExpandAll` function.

**Demonstrating orthogonality.** We integrate the product of two Legendre polynomials of different degree. Consider

$$\int_{-1}^1 P_2(x)P_4(x) dx$$

Substituting explicit polynomials for the  $P_n(x)$  gives

$$\int_{-1}^1 \left( \frac{-1 + 3x^2}{2} \right) \left( \frac{3 - 30x^2 + 35x^4}{8} \right) dx = \int_{-1}^1 \left( -\frac{3}{16} + \frac{39x^2}{16} - \frac{125x^4}{16} + \frac{105x^6}{16} \right) dx$$

Then term by term integration gives

$$-\frac{3}{8} + \frac{13}{8} - \frac{25}{8} + \frac{15}{8} = 0$$

5. `beginRunonGroup`, `beginLeftAlignedGroup`, and `endRunonGroup`, `endLeftAlignedGroup` produce analogous effects.
6. `integrateTermByTermAndHold` distributes integration over the `Plus` in the integrand, and stops the result coalescing.
7. `allowEvaluation` removes the `Hold` around each term.

A large variety of worked examples to help teach orthogonality of special functions can be generated by applying `autorecord` to input files that follow the general style of this prototype, use the same `biLo` functions, and cycle through sets of parameters.

## 9. Directives

The control file directives are simply statements that use raw MATHEMATICA or invoke procedures

in our `bi`, `forTeX` and related packages. Thus, in Section 2, `bi` sorting functions are assigned to `format`—a surrogate head that is applied to each output expression before it is encoded. The directives in Section 8 are the names of  $\TeX$  functions that put the delimiters for display math into the output, and control several program switches.

Statements are written as directives that:

1. substitute  $\TeX$ -like names for other identifiers,
2. convert function expressions that carry indexes as ordinary arguments into representations that lead to sub- and superscripting,
3. change font and impose bracketing,
4. fine tune the style as described in Section 5,
5. perform substantive steps that are not typeset.

There is further flexibility. `autorecord` is used in the `displayBoth` mode to document the symbolic algebra part of a mechanized derivation. It plays back the input in typewriter font and the evaluated expressions as displayed mathematics. Options allow typewriter font for both. Each input expression can be run-on or left-aligned above the “ $\sim$ ” symbol and conventionally displayed result. Further styles of alignment are provided for the `outputOnly` and `outputBoth` modes. Provision can be made to number the equations.

Options can be changed dynamically. Also, the directives can be made conditional on flags that are set at execution time, to vary the style of output from different runs using the same control file. The content can be varied by conditional directives that make evaluation occur silently until balancing directives are reached, or bypass text and/or evaluation completely. This can be used, e.g. to produce terse and detailed derivations from a single file, or problem sets with and without solutions.

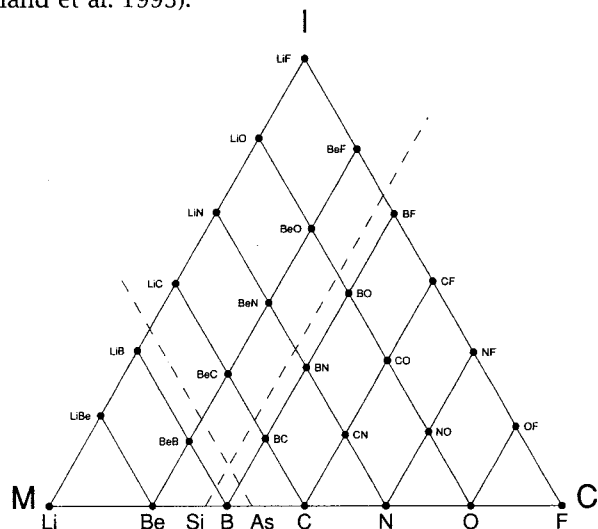
Four pre- and postprocessing functions, that default to `Identity`, which returns the argument unchanged, provide part of the flexibility.

1. `ipc`—this converts input statements into a preferred form that is recorded,
2. `ivc`—this converts input statements from a variant notation into valid `MATHEMATICA`,
3. `opc`—this acts on the immediate result of `MATHEMATICA` evaluation to give output of preferred appearance,
4. `tpc`—this operates on each block of text before it is displayed.

These parallel, in part the built-in `$PreRead`, `$Pre`, `$Post` and `$PrePrint` functions. `autorun` can document the full extent of activity by displaying any combination of the input as read, the restyled input, the input to `MATHEMATICA`, the output from `MATHEMATICA` and the restyled output, alongside tags that default to “In”, “in”, “IN”, “Out” and “out”, respectively.

## 10. Using graphics

`MATHEMATICA` has extensive graphics capabilities. It saves graphics objects as `POSTSCRIPT` files. `psfig` imports these to  $\TeX$  documents and is very easy to use. The “van Arkel – Ketelaar” diagram below relates to binary compounds of the chemical elements and an atomic property called “electronegativity” (Allen, Leland et al. 1993).



I/M/C character of 1st row binaries

The letters I, M and C at the corners of the triangle stand for ionic, metallic and covalent. We wrote a `MATHEMATICA` procedure that reads a list of elements and produces the corresponding diagram. This example involves the algorithmic placement of text, with provision for interactive fine tuning.

The ability to import graphics helps check symbolic calculations. Computer algebra is fraught with possible error, though this is seldom discussed. The early stages of a symbolic calculation often produce results that are in the literature, and there is merit to making a comparison.

Our first efforts to check results this way showed that, often, in hand generated formulas, a particular subexpression is represented by a single variable in one place but not another. Machine generated formulas tend to be more consistent. `bi` was developed, in part, to target the changes that are needed to bring the content and arrangement of computed results into agreement with published material.

Even when this has been done, however, the visual comparison can be tedious. For example, as part of a longer calculation (Barnett 1994a), we reconstructed a 33-term recurrence formula, that occupies half a page of *Physical Review* (Pekeris 1958). It can be written:

$$\sum_{(\lambda, \mu, \nu) \in S} c_{\lambda, \mu, \nu}(l, m, n) A(l + \lambda, m + \mu, n + \nu) = 0$$

where the  $c$ 's are simple algebraic expressions in  $(l, m, n)$ , and  $S$  consists of 33 triples whose elements are between  $-2$  and  $2$ . To help other workers see that the results are in agreement, the relevant passage in Physical Review was scanned, and the image dissected into about 40 pieces that contain separate terms or, in some cases, parts of terms, using a suite of image processing tools written locally. The pieces were imported into a document above the respective terms generated by MATHEMATICA, using another utility. This has been done here for terms 1-3 and 32, 33, by conversion to POSTSCRIPT and use of psfig/tex. Scaling is varied slightly.

$$\begin{aligned}
 &4(l+1)(l+2)[-Z+\epsilon(1+m+n)]A(l+2, m, n) \\
 &4(l+1)(l+2)[-Z+\epsilon(1+m+n)]A(l+2, m, n) \\
 &+4(m+1)(m+2)[-Z+\epsilon(1+l+n)]A(l, m+2, n) \\
 &4(m+1)(m+2)[-Z+\epsilon(1+l+n)]A(l, m+2, n) \\
 &+4(l+1)(m+1)[1-2Z+\epsilon(2+l+m)]A(l+1, m+1, n) \\
 &4(l+1)(m+1)[1-2Z+\epsilon(2+l+m)]A(l+1, m+1, n) \\
 &\dots \\
 &+2ln[1-2Z+\epsilon(2m+n+1)]A(l-1, m, n-1) \\
 &2ln[1-2Z+\epsilon(2m+n+1)]A(l-1, m, n-1) \\
 &+2mn[1-2Z+\epsilon(2l+n+1)]A(l, m-1, n-1) \\
 &2mn[1-2Z+\epsilon(2l+n+1)]A(l, m-1, n-1)
 \end{aligned}$$

This technique can be used to publish scanned pictures together with algorithmically constructed graphics in many areas of work.

Setting the output of symbolic computation, using special fonts and macros for chemical formulas, chess situations, musical scores and other iconic displays has still further potential. The use of built-up formulas by organic and biological chemists is well known. These, and the further needs of inorganic and solid state chemists are discussed in (Jensen 1989). We have made a start on interfacing MATHEMATICA with the work reported earlier (Haas and O'Kane 1987; Tutelaers 1992; Taupin 1992).

As a simple example of non-mathematical output produced by MATHEMATICA, two very short procedures that construct a word count and encode it as a table produced the following display from a 10-line passage in Macbeth.

a	3	adder's	1	and	8
bake	1	bat	1	blind	1
...					
trouble	2	wing	1	wool	1
worm's	1				

## Distribution

The `bi10` package is in, and the `forTeX` material soon will be in the anonymous ftp library at:

mondrian.princeton.edu (128.112.224.14)

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